

# Repunit Lehmer numbers

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## Abstract

A Lehmer number is a composite positive integer  $n$  such that  $\phi(n) \mid n - 1$ . In this paper, we show that given a positive integer  $g > 1$  there are at most finitely many Lehmer numbers which are repunits in base  $g$ , and they are all effectively computable. Our method is effective and we illustrate it by showing that there is no such Lehmer number when  $g \in [2, 1000]$ .

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## 1 Introduction

Let  $\phi(n)$  be the Euler function of the positive integer  $n$ . Clearly,  $\phi(n) = n - 1$  if  $n$  is a prime. Lehmer [4] (see also B37 in [3]) conjectured that if  $\phi(n) \mid n - 1$ , then  $n$  is prime. To this day, no counterexample to this conjecture has been found. A composite number  $m$  such that  $\phi(m) \mid m - 1$  is called a *Lehmer number*. Thus, Lehmer's conjecture is that Lehmer numbers don't exist but it is not even known that there should be at most finitely many of them.

Given a positive integer  $g > 1$  a base  $g$  repunit is a number of the form  $m = (g^n - 1)/(g - 1)$  for some integer  $n \geq 1$ . We will refer to such numbers

simply as repunits without mentioning the dependence on  $g$ . It is not known whether given  $g$  there are infinitely many repunit primes. When  $g = 2$  such primes are better known as Mersenne primes. In [5], it was shown that there is no Lehmer number in the Fibonacci sequence. Here, we use some ideas from [5] together with finer arguments to prove the following results. In what follows, we write  $u_n = (g^n - 1)/(g - 1)$ .

**Theorem 1.** *For each fixed  $g > 1$ , there are only finitely many positive integers  $n$  such that  $u_n$  is a Lehmer number, and all are effectively computable.*

**Theorem 2.** *There is no Lehmer number of the form  $u_n$  when  $2 \leq g \leq 1000$ .*

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## 2 Preliminaries

For a prime  $q$  and a nonzero integer  $m$  we write  $\nu_q(m)$  for the exponent of  $q$  in the factorization of  $m$ . We start by collecting some elementary and well-known properties of the sequence of general term  $u_n = (g^n - 1)/(g - 1)$  for  $n \geq 1$ .

**Lemma 1.** *i)  $u_n = g^{n-1} + \dots + g + 1$ . In particular,  $u_n$  is coprime to  $g$ .*

*ii) The sequence  $u_n$  satisfies the linear recurrence*

$$u_1 = 1, \quad u_n = gu_{n-1} + 1, \quad n \geq 2. \quad (1)$$

*iii) If  $d \mid n$ , then  $u_d \mid u_n$ .*

*iv) Let  $q$  be a prime. If  $q \mid n$ , then  $q \mid \phi(u_n)$ .*

*v) Let  $q$  be a prime not dividing  $g$ . If  $q \mid n$ , then  $\nu_q(u_{n-1}) \leq \nu_q(u_f) \leq \nu_q(u_{q-1})$ , where  $f$  is the order of  $g$  modulo  $q$ .*

*vi) If  $u_n$  is a Lehmer number, then  $(u_n, g - 1) = 1$ .*

*Proof.* *i)* and *ii)* are obvious. For *iii)*, we observe that

$$u_n = \frac{g^n - 1}{g - 1} = \frac{(g^d)^{n/d} - 1}{g^d - 1} \cdot \frac{g^d - 1}{g - 1} = \left( (g^d)^{\frac{n}{d}-1} + \dots + 1 \right) u_d.$$

*iv)* If  $q = 2$ , then  $u_n \geq u_2 = g + 1 > 2$ , therefore  $\phi(u_n)$  is even. Assume now that  $q$  is odd. Let  $p$  be a prime which divides  $u_q$ . Then,  $g^q \equiv 1 \pmod{p}$ , so the order of  $g$  modulo  $p$  is 1 or  $q$ . If it is  $q$ , then  $q \mid p - 1 \mid \phi(u_q)$ . Since by *iii)* we know that  $u_q \mid u_n$ , we get that  $q \mid \phi(u_q) \mid \phi(u_n)$ , which is what we wanted. Assume now that the order of  $g$  modulo  $p$  is 1 for all primes  $p$  dividing  $u_q$ . Let us show that this cannot happen. If it would, then  $p \mid g - 1$  for all such primes  $p$ . Since also  $p \mid u_q$ , we have

$$0 \equiv u_q \equiv \frac{g^q - 1}{g - 1} = g^{q-1} + \dots + g + 1 \equiv 1 + \dots + 1 + 1 \equiv q,$$

where all congruences above are modulo  $p$ . Thus,  $p \mid q$ , therefore  $p = q$ . Hence,  $u_q = q^\alpha$  for some positive integer  $\alpha$ . However, writing  $g - 1 = q\lambda$  with some positive integer  $\lambda$ , we get

$$\begin{aligned} u_q &= (1 + q\lambda)^{q-1} + (1 + q\lambda)^{q-2} + \dots + (1 + q\lambda) + 1 \\ &\equiv (1 + (q-1)q\lambda) + (1 + (q-2)q\lambda) + \dots + (1 + q\lambda) + 1 \pmod{q^2} \\ &\equiv q + q\lambda((q-1) + \dots + 1) \pmod{q^2} \\ &\equiv q + \frac{q^2(q-1)\lambda}{2} \pmod{q^2} \\ &\equiv q \pmod{q^2}. \end{aligned}$$

In the above chain of congruences, we used the fact that  $q$  is odd, therefore  $(q-1)/2$  is an integer. The above argument shows that  $q \parallel u_q$ ; hence,  $\alpha = 1$ . So,  $u_q = q$ . However, we clearly have  $u_q \geq 2^q - 1 > q$ , which is a contradiction.

*v)* We may also assume that  $q \mid u_{n-1}$ , otherwise  $\nu_q(u_{n-1}) = 0$  and the first inequality is clear. Now  $g^{n-1} \equiv 1 \pmod{q}$ , and so  $f \mid n - 1$ . We now write

$$u_{n-1} = \left( (g^f)^{\frac{n-1}{f}-1} + \dots + 1 \right) u_f.$$

The quantity in brackets above is not divisible by  $q$  since it is congruent to  $(n-1)/f$  modulo  $q$  and  $q \mid n$ . Thus,  $\nu_q(u_{n-1}) \leq \nu_q(u_f) \leq \nu_q(u_{q-1})$ , where the last inequality follows because  $f \mid q - 1$ , so,  $u_f \mid u_{q-1}$  by *iii)*.

*vi)* Suppose that  $q$  is a prime dividing both  $u_n$  and  $g - 1$ . We then have that  $g \equiv 1 \pmod{q}$  and  $u_n = g^{n-1} + \dots + 1 \equiv n \pmod{q}$ . Thus,  $q \mid n$ . By

*iv*), we know that  $q \mid \phi(u_n)$ . Since  $u_n$  is a Lehmer number, we know that  $\phi(u_n) \mid u_n - 1 = gu_{n-1}$ . Since  $q$  divides  $g - 1$ , it cannot divide  $g$ , therefore  $q \mid u_{n-1}$ . Hence,  $q \mid u_n - u_{n-1} = g^{n-1}$ , which is not possible.  $\square$

In the next lemma, we gather some known facts about Lehmer numbers.

**Lemma 2.** *i) Any Lehmer number must be odd and square-free.*

*ii) If  $m = p_1 \cdots p_K$  is a Lehmer number, then  $K^{2^K} > m$ .*

*iii) If  $m = p_1 \cdots p_K$  is a Lehmer number, then  $K \geq 14$ .*

*Proof.* *i)* If  $m > 2$  then  $\phi(m)$  is even, and since  $\phi(m) \mid m - 1$ , we get that  $m$  must be odd. If  $p^2 \mid m$ , then  $p \mid \phi(m)$ , and since  $\phi(m) \mid m - 1$ , we have  $p \mid m - 1$ , which is not possible. Part *ii)* was proved by Pomerance in [6], while part *iii)* was proved by Cohen and Hagis in [2].  $\square$

**Lemma 3.** *Theorems 1 and 2 hold when  $g$  is even.*

*Proof.* Note that

$$2^K \mid (p_1 - 1) \cdots (p_K - 1) = \phi(u_n) \mid u_n - 1 = gu_{n-1}.$$

We observe that if  $g$  is even, then  $u_{n-1}$  is odd. In that case, we have

$$K \leq \nu_2(\phi(u_n)) \leq \nu_2(gu_{n-1}) = \nu_2(g), \quad (2)$$

implying, by Lemma 2 *ii)*, that

$$g^{n-1} < u_n < K^{2^K} \leq (\nu_2(g))^{2^{\nu_2(g)}} \leq (\nu_2(g))^g.$$

Thus,

$$n \leq 1 + \left\lfloor \frac{g \log(\nu_2(g))}{\log g} \right\rfloor.$$

For Theorem 2, we observe that  $\nu_2(g) \leq 9$  for any  $g \leq 1000$ , and we obtain a contradiction from (2) and Lemma 2 *iii)*.  $\square$

From Lemma 1 *i)*, we see that if  $g$  is odd and  $n$  is even, then  $u_n$  is even, so Lemma 2 *i)* shows that  $u_n$  cannot be a Lehmer number. From now on, we shall assume that both  $g$  and  $n$  are odd and larger than 1 and that  $u_n = (g^n - 1)/(g - 1)$  is a Lehmer number. We also keep the notation:

$$n = q_1^{\alpha_1} \cdots q_s^{\alpha_s}, \quad \text{where } 2 < q_1 < \cdots < q_s \quad (3)$$

are primes and  $\alpha_1, \dots, \alpha_s$  are positive integers, and

$$u_n = p_1 \cdots p_K, \quad \text{where } 2 < p_1 < \cdots < p_K \quad (4)$$

are also primes.

### 3 Proof of Theorem 1

#### 3.1 Primitive divisors

Let  $(A_n)_{n \geq 1}$  denote a sequence with integer terms. We say that a prime  $p$  is a *primitive divisor* of  $A_n$  if  $p \mid A_n$  and  $\gcd(p, A_m) = 1$  for all non-zero terms  $A_m$  with  $1 \leq m < n$ .

In 1886, Bang [1] showed that if  $g > 1$  is any fixed integer, then the sequence  $(A_n)_{n \geq 1}$  of  $n$ th term  $A_n = g^n - 1$  has a primitive divisor for any index  $n > 6$ .

We will apply this important theorem to our sequence  $u_n$ .

**Lemma 4.** *If  $d > 1$  is odd, then  $u_d$  has a primitive divisor  $p_d$ . Furthermore,  $p_d \equiv 1 \pmod{2d}$ .*

*Proof.* We revisit the argument already used at Lemma 1 *iv*). We write  $v_n = g^n - 1$ . It is well-known that  $\gcd(v_n, v_m) = v_{\gcd(n, m)}$ . Observe also that

$$\frac{v_d}{v_1} = u_d = g^{d-1} + \cdots + 1 \equiv d \pmod{g-1},$$

therefore if  $d$  is a prime not dividing  $g-1$ , then  $v_d$  has primitive divisors. If  $d > 2$  is a prime dividing  $g-1$ , then the above argument, or the argument from the proof of Lemma 1 *iv*), shows that  $\gcd(v_d, v_1)$  is a power of  $d$ . Write  $g-1 = d\lambda$  and observe that

$$\begin{aligned} \frac{v_d}{v_1} &= (1 + d\lambda)^{d-1} + (1 + d\lambda)^{d-2} + \cdots + 1 \\ &\equiv (1 + (d-1)d\lambda) + (1 + (d-2)d\lambda) + \cdots + 1 \\ &= d + d\lambda((d-1) + (d-2) + \cdots + 1) \pmod{d^2} \\ &\equiv d + \frac{d^2(d-1)}{2}\lambda \pmod{d^2} \equiv d \pmod{d^2}. \end{aligned}$$

Thus,  $d \parallel v_d/v_1$ , and therefore

$$\frac{v_d}{dv_1} = \frac{1}{d}(g^{d-1} + \cdots + 1) > 1$$

is an integer coprime to  $v_1$ , so  $v_d$  again has primitive divisors. Thus,  $v_3$  and  $v_5$  (and, of course,  $v_1$  if  $g > 2$ ) have primitive divisors. The fact that  $v_d$  has primitive divisors for all odd  $d \geq 7$  follows from Bang's result.

We now note that if  $p$  is a primitive prime divisor of  $v_d$  for  $d > 1$ , then  $g^d \equiv 1 \pmod{p}$ , and  $d$  is the order of  $g \pmod{p}$ . Indeed, for if not, then

$f < d$  and  $p \mid v_f$ , contradicting the fact that  $p$  is primitive for  $v_d$ . So,  $d \mid p-1$ , and since  $d$  is odd, we get that  $d \mid (p-1)/2$ . Thus,  $p \equiv 1 \pmod{2d}$ .

Since a prime factor of  $g-1$  cannot be a primitive divisor for  $v_d$  except for  $d=1$ , we deduce that if  $d > 1$ , then the primitive prime divisors for  $v_d$  are exactly those of  $u_d = v_d/(g-1)$ , and we get the first assertion of the lemma.  $\square$

In what follows, for a positive integer  $m$  we use  $\omega(m)$  and  $\tau(m)$  for the number of prime divisors and the total number of divisors of  $m$ , respectively.

**Lemma 5.** *If  $u_n$  is square-free,  $n$  is odd and  $(u_n, g-1) = 1$ , then*

$$\begin{aligned} \log \left( \frac{u_n}{\phi(u_n)} \right) &< \frac{\omega(n)}{2q} \left( 1 + \log \left( \frac{q \log g}{\log(2q+1)} \right) \right) \\ &+ \frac{\tau(n) - 2}{2q^2} \left( 1 + \log \left( \frac{q^2 \log g}{\log(2q^2+1)} \right) \right), \end{aligned}$$

where  $q$  is the smallest prime dividing  $n$ .

*Proof.* We write  $\mathcal{P}_d = \{p \text{ is primitive prime divisor for } u_d\}$ . We shall first prove that

$$\prod := \prod_{1 < d \mid n} \prod_{p \in \mathcal{P}_d} p = u_n.$$

To see the above formula, we observe that if  $p \mid u_d$  and  $p \nmid g-1$ , then  $p \in \mathcal{P}_d$  for some  $1 < d \mid n$ . Since  $u_n$  is square-free, we have that  $u_n \mid \prod$ . On the other hand, the sets  $\mathcal{P}_d$  are disjoint, and if  $p \in \mathcal{P}_d$ , then  $p \mid u_d \mid u_n$ . Thus,  $\prod \mid u_n$ .

Now, since  $u_n$  is square-free,

$$\phi(u_n) = \prod_{1 < d \mid n} \prod_{p \in \mathcal{P}_d} (p-1),$$

and then

$$\log \left( \frac{u_n}{\phi(u_n)} \right) < \sum_{\substack{d \mid n \\ d > 1}} \sum_{p \in \mathcal{P}_d} \frac{1}{p-1}.$$

Since all the primes  $p \in \mathcal{P}_d$  are congruent to 1 (mod  $2d$ ), we have

$$S_d := \sum_{p \in \mathcal{P}_d} \frac{1}{p-1} \leq \frac{1}{2d} \sum_{j=1}^{\#\mathcal{P}_d} \frac{1}{j} \leq \frac{1}{2d} (1 + \log \#\mathcal{P}_d).$$

To bound the cardinality of  $\mathcal{P}_d$ , we observe that  $(2d+1)^{\#\mathcal{P}_d} \leq u_d < g^d$ , so

$$\#\mathcal{P}_d < \frac{d \log g}{\log(2d+1)}.$$

We observe that  $d \geq q$  and if  $d$  is not a prime, then  $d \geq q^2$ . Then

$$\begin{aligned} \sum_{1 < d|n} S_d &= \sum_{\substack{d|n \\ d \text{ prime}}} S_d + \sum_{\substack{d|n \\ d \text{ composite}}} S_d \leq \omega(n) \frac{1}{2q} \left( 1 + \log \left( \frac{q \log g}{\log(2q+1)} \right) \right) \\ &\quad + (\tau(n) - 2) \frac{1}{2q^2} \left( 1 + \log \left( \frac{q^2 \log g}{\log(2q^2+1)} \right) \right). \end{aligned}$$

□

### 3.2 Bounds for $q_1$ and $\tau(n)$

Recall that we keep the notations from (3) and (4).

**Lemma 6.** *If  $u_n$  is a Lehmer number and  $n$  is odd, then*

$$\begin{aligned} \tau(n/q_i) &\leq \frac{\alpha_i(\alpha_i+1)}{2} \tau(n/q_i^{\alpha_i}) \leq \nu_{q_i}(\phi(u_n)) \leq \nu_{q_i}(gu_{n-1}) \\ &\leq \begin{cases} \nu_{q_i}(g), & \text{if } q_i | g; \\ \nu_{q_i}(u_{q_i-1}), & \text{if } q_i \nmid g \end{cases} \end{aligned} \quad (5)$$

for all  $i = 1, \dots, s$ .

*Proof.* Lemma 4 implies that for each divisor of  $n$  of the form  $q_i^\alpha d$  with  $1 \leq \alpha \leq \alpha_i$  and  $d | (n/q_i^{\alpha_i})$ , the divisor  $u_{q_i^\alpha d}$  of  $u_n$  has a primitive prime factor  $p_{q_i^\alpha d} \equiv 1 \pmod{dq_i^\alpha}$ . In particular,  $q_i^\alpha | p_{dq_i^\alpha} - 1$ , and the primes  $p_{dq_i^\alpha}$  are distinct as  $d$  ranges over the divisors of  $n/q_i^{\alpha_i}$ . Thus,

$$\begin{aligned} q_i^{(1+\dots+\alpha_i)\tau(n/q_i^{\alpha_i})} &| \prod_{1 \leq \alpha \leq \alpha_i} \prod_{d|n/q_i^{\alpha_i}} (p_{dq_i^\alpha} - 1) | \prod_{p|u_n} (p-1) \\ &= \phi(u_n) | u_n - 1 | gu_{n-1}, \end{aligned}$$

which gives the two central inequalities. The first inequality is trivial and the equality holds when  $\alpha_i = 1$ . When  $q_i | g$ , the last inequality follows from Lemma 1 *i*), while when  $q_i \nmid g$ , then  $\nu_{q_i}(gu_{n-1}) = \nu_{q_i}(u_{n-1})$ , and we apply Lemma 1 *v*) to get the desired conclusion. □

**Lemma 7.** *Let  $u_n$  be a Lehmer number with both  $n$  and  $g$  odd. If  $q_i > \sqrt{g}$ , then*

$$\tau(n/q_i) \leq q_i - 2.$$

*Proof.* If  $q_i \mid g$  and  $q_i > \sqrt{g}$ , then  $\nu_{q_i}(g) = 1$ , and Lemma 6 above gives

$$\tau(n/q_i) \leq \nu_{q_i}(g) = 1 \leq q_i - 2. \quad (6)$$

If  $q_i \nmid g$ , then, again by Lemma 6 above, we have

$$\tau(n/q_i) \leq \nu_{q_i}(u_{q_i-1}).$$

Observe that

$$u_{q_i-1} \mid g^{q_i-1} - 1 = \left(g^{(q_i-1)/2} - 1\right) \left(g^{(q_i-1)/2} + 1\right).$$

Since  $q_i$  cannot divide both factors above, we have that

$$\tau(n/q_i) \leq \nu_{q_i}(g^{(q_i-1)/2} + \epsilon) \quad \text{for some } \epsilon \in \{-1, +1\}.$$

If  $\tau(n/q_i) \geq q_i - 1$ , then

$$q_i^{q_i-1} \leq q_i^{\tau(n/q_i)} \leq g^{(q_i-1)/2} + 1 \leq (q_i^2 - 1)^{(q_i-1)/2} + 1, \quad (7)$$

and we get a contradiction for  $q_i > 3$ , because

$$q_i^{q_i-1} = ((q_i^2 - 1) + 1)^{(q_i-1)/2}$$

and the expression on the right is larger than  $(q_i^2 - 1)^{(q_i-1)/2} + 1$  except when  $q_i = 3$ .

Finally, if  $q_i = 3$ , the only odd  $g < q_i^2$  with  $q_i \nmid g$  are  $g = 5$  and  $g = 7$ . But in both cases we have  $\tau(n/3) \leq \nu_3(u_2) \leq 1 \leq q_i - 2$ , which completes the proof of this lemma.  $\square$

**Lemma 8.** *Let  $u_n$  be a Lehmer number with both  $n$  and  $g$  odd. Then*

$$q_1 \leq \max\{\sqrt{g}, 19\}. \quad (8)$$

*Proof.* Assume that the above inequality does not hold. Then  $q_1 \geq 23$ ,  $g \leq q_1^2 - 1$ , and since  $q_1 > \sqrt{g}$ , we can apply Lemma 7 to deduce that  $\tau(n) \leq 2\tau(n/q_1) \leq 2q_1 - 4$ . We also observe that  $\tau(n) \geq 2^{\omega(n)}$ , so  $\omega(n) \leq \log(2q_1 - 4)/\log 2$ .



Since  $u_n$  is a Lehmer number, we have that  $2 \leq u_n/\phi(u_n)$ . Now Lemma 5 and the bounds above give

$$\begin{aligned} \log 2 &< \frac{\log((2q_1 - 4)/\log 2)}{2q_1} \left( 1 + \log \left( \frac{q_1 \log(q_1^2 - 1)}{\log(2q_1 + 1)} \right) \right) \\ &+ \frac{2q_1 - 6}{2q_1^2} \left( 1 + \log \left( \frac{q_1^2 \log(q_1^2 - 1)}{\log(2q_1^2 + 1)} \right) \right), \end{aligned}$$

which is false for  $q_1 \geq 23$ .  $\square$

For a given value of  $g$ , Lemma 8 gives us our bound for  $q_1$  and then this is used in Lemma 6, since  $\tau(n) \leq 2\tau(n/q_1)$ , to give a bound for  $\tau(n)$ . Observe also that  $\omega(n) \leq \log \tau(n)/\log 2$ .

### 3.3 The conclusion of the proof of Theorem 1

Since we have already proved that both  $s = \omega(n)$  and  $\tau(n)$  are bounded by effectively computable constants depending only on  $g$ , in order to conclude the proof of Theorem 1 it is enough to prove that all the primes  $q_i$  with  $i = 1, \dots, s$  are also bounded by effectively computable constants depending on  $g$ . We shall prove this by induction on  $i = 1, \dots, s$  observing that this has already been achieved for  $i = 1$ . Let  $i \leq s - 1$  and assume that  $q_i$  has been bounded. Put  $Q_i = \prod_{j=1}^{j=i} q_j^{\alpha_j}$ . There are only finitely many possibilities for this number. We put  $g_i = g^{Q_i}$ ,  $n_i = n/Q_i$  and rewrite the condition that  $u_n$  is Lehmer as

$$a\phi \left( \frac{g^{Q_i} - 1}{g - 1} \cdot \frac{g_i^{n_i} - 1}{g_i - 1} \right) = u_n - 1 = \frac{g^{Q_i} - 1}{g - 1} \cdot \frac{g_i^{n_i} - 1}{g_i - 1} - 1$$

with some integer  $a \geq 2$ . We put  $w_m = (g_i^m - 1)/(g_i - 1)$  for the sequence of repunits in base  $g_i$ . Then, since  $u_n$  is square-free, we get that

$$a\phi(u_{Q_i})\phi(w_{n_i}) = u_{Q_i}w_{n_i} - 1,$$

therefore

$$a \frac{\phi(u_{Q_i})}{u_{Q_i}} = \frac{w_{n_i}}{\phi(w_{n_i})} - \frac{1}{u_{Q_i}\phi(w_{n_i})}. \quad (9)$$

The left hand side takes only finitely many values, which are all effectively computable. Assume that it takes some value  $\delta \leq 1$ . Then

$$w_{n_i} - 1 < w_{n_i} - \frac{1}{u_{Q_i}} = \delta\phi(w_{n_i}) \leq \phi(w_{n_i}),$$

a contradiction. Thus, it remains to study the case when the right hand side in (9) is  $> 1$ . Let  $\delta_i > 1$  be the smallest possible value larger than 1 of the left hand side of (9). Clearly, this is effectively computable. We then get

$$\delta_i < \frac{w_{n_i}}{\phi(w_{n_i})}.$$

We observe that  $w_{n_i}$  is a sequence “like”  $u_n$  but the new value of  $g$  is  $g_i = g^{Q_i}$  and the new value of  $n$  is  $n_i = n/Q_i$ . Thus, the smallest prime factor of  $n_i$  is  $q_{i+1}$ . We also note that  $\tau(n_i) = \tau(n/Q_i) < \tau(n)$  which is bounded, and that  $\omega(n_i) < \omega(n)$ . Finally, we observe that  $(w_{n_i}, g^{Q_i} - 1) = 1$ , otherwise, since  $(w_{n_i}, g - 1) = 1$ , the number  $u_n = (g^{Q_i} - 1)w_{n_i}/(g - 1)$  would not be square-free.

We now apply Lemma 5 to obtain that

$$\begin{aligned} \log \delta_i &< \frac{\omega(n_i)}{2q_{i+1}} \left( 1 + \log \left( \frac{Q_i q_{i+1} \log g}{\log(2q_{i+1} + 1)} \right) \right) \\ &+ \frac{\tau(n_i) - 2}{2q_{i+1}^2} \left( 1 + \log \left( \frac{Q_i q_{i+1}^2 \log g}{\log(2q_{i+1}^2 + 1)} \right) \right). \end{aligned} \quad (10)$$

Hence,  $\log \delta_i \ll \frac{\log q_{i+1}}{q_{i+1}}$ , where the constant implied by the Vinogradov symbol  $\ll$  above depends only on  $g$ , implying that  $q_{i+1}$  must be bounded by some effectively computable constant depending only on  $g$ . This concludes the proof of Theorem 1.

## 4 Proof of Theorem 2

We assume that  $g$  is odd and  $3 \leq g \leq 999$ , so that  $3 \leq q_1 \leq 31$  by Lemma 8.

**Claim 1:** That  $\nu_{q_1}(u_{q_1-1}) \leq 5$  can be checked with Mathematica. In particular, by Lemma 6, we have that if  $q_1 \nmid g$ , then  $\nu_{q_1}(\phi(u_n)) \leq 5$ .

**Claim 2:**  $\tau(n/q_1) \leq \nu_{q_1}(\phi(u_n)) \leq 6$ , and  $s \leq 3$ .

Suppose first that  $q_1 \mid g$ . Then, by Lemma 6,

$$\tau(n/q_1) \leq \nu_{q_1}(\phi(u_n)) \leq \nu_{q_1}(gu_{n-1}) = \nu_{q_1}(g) \leq \left\lfloor \frac{\log g}{\log q_1} \right\rfloor \leq \left\lfloor \frac{\log 1000}{\log 3} \right\rfloor = 6.$$

In the above, in fact  $\nu_{q_1}(g) < 6$  unless  $(q_1, g) = (3, 729)$ . Then, for any  $q_1$ , by Claim 1, either  $q_1 = 3$  and  $\tau(n/q_1) \leq 6$ , or  $\tau(n/q_1) \leq 5$ . In particular,  $\tau(n) \leq 2\tau(n/q_1) \leq 12$ , which shows that  $s \leq 3$ .

**Claim 3:**  $s \geq 2$ .

Let us see indeed that for our particular case we cannot have  $s = 1$ . If this were so, then  $n = q_1^{\alpha_1}$ . Then each prime factor  $p_j$  of  $u_n$  is primitive for some divisor  $d > 1$  of  $n$ , which is a power of  $q_1$  (again, this is because  $\gcd(u_n, g - 1) = 1$ ). Thus,  $p_j \equiv 1 \pmod{q_1}$  for all  $j = 1, \dots, K$ , showing that  $\nu_{q_1}(\phi(u_n)) \geq K \geq 14$  (see Lemma 2 *iii*), which contradicts the fact that  $\nu_{q_1}(\phi(u_n)) \leq 6$ . Hence,  $s \geq 2$ .

**Claim 4:**  $\alpha_1 = 1$  except when  $(\alpha_1, q_1, g) = (2, 3, 729)$ .

Put again, as in the proof of Theorem 1,  $Q_1 = q_1^{\alpha_1}$ . By Lemma 6 and the fact that  $s \geq 2$ , we have

$$\alpha_1(\alpha_1 + 1) \leq \frac{\alpha_1(\alpha_1 + 1)}{2} \tau(n/q_1^{\alpha_1}) \leq \nu_{q_1}(\phi(u_n)).$$

By Claims 1 and 2 above, we know that  $\nu_{q_1}(\phi(u_n)) \leq 5$ , except when  $(\alpha_1, q_1, g) = (2, 3, 729)$ . So,  $\alpha_1 = 1$  except for this case.

Note that, at any rate, since  $s \geq 2$ , it follows that  $2 \leq \tau(n/q_1) \leq \nu_{q_1}(gu_{q_1-1})$ . A computation with Mathematica revealed 431 possibilities for the pairs  $(q_1, g)$  in our range satisfying  $\nu_{q_1}(gu_{q_1-1}) \geq 2$ .

**Claim 5:**  $q_2 \leq 19$ .

The smallest left hand side in (9) computed over all the 432 possible pairs  $(Q_1, g)$  has  $\delta_1 > 1.49$  (it was obtained for  $g = 809$ ,  $Q_1 = q_1 = 3$  and  $a = 2$ , for which the obtained value is  $> 1.495$ ). Of course, we did not factor all the numbers of the form  $(g^{Q_1} - 1)/(g - 1)$ . If  $q_1 = 31$ , then the smallest prime  $p_1 \equiv 1 \pmod{q_1}$  is 311. The number  $K$  of prime factors of  $u_{31}$  satisfies therefore

$$K < \frac{\log u_{q_1}}{\log p_1} < \frac{3 \cdot 31 \cdot \log 10}{\log 311} < 38;$$

hence,

$$a \frac{\phi(u_{q_1})}{u_{q_1}} \geq 2 \left(1 - \frac{1}{311}\right)^{37} > 1.7.$$

Similarly, using the fact that when  $q_1 = 29$  and 23 the first two primes congruent to 1  $\pmod{q_1}$  are 59 and 233, and 47 and 139 respectively, and

$$\frac{3 \cdot 29 \cdot \log 10}{\log 233} < 37 \quad \text{and} \quad \frac{3 \cdot 23 \cdot \log 10}{\log 139} < 33,$$

we have that

$$\begin{aligned} a \frac{\phi(u_{q_1})}{u_{q_1}} &\geq 2 \min \left\{ \left(1 - \frac{1}{59}\right) \left(1 - \frac{1}{233}\right)^{36}, \left(1 - \frac{1}{47}\right) \left(1 - \frac{1}{139}\right)^{32} \right\} \\ &> 1.55, \end{aligned}$$

whenever  $q_1 \in \{23, 29\}$ . Thus, we have factored only the numbers  $u_{Q_1}$  with  $Q_1 \leq 19$ . We now use inequality (10) for  $i = 1$  to obtain

$$\begin{aligned} \log(1.49) &< \frac{\omega(n_1)}{2q_2} \left(1 + \log \left(\frac{Q_1 q_2 \log g}{\log(2q_2 + 1)}\right)\right) \\ &\quad + \frac{\tau(n_1) - 2}{2q_2^2} \left(1 + \log \left(\frac{Q_1 q_2^2 \log g}{\log(2q_2^2 + 1)}\right)\right). \end{aligned}$$

If  $q_1 > 3$ , then  $Q_1 = q_1 \leq 31$ . If  $q_1 = 3$ , then  $Q_1 = q_1^2 = 9$ . Thus,  $Q_1 \leq 31$  in both cases. We also saw in Claims 1 and 2 that  $\tau(n_1) \leq \tau(n/q_1) \leq 6$ , so also  $\omega(n_1) \leq 2$ . Hence,

$$\log(1.49) < \frac{1}{q_2} \left(1 + \log \left(\frac{31q_2 \log 999}{\log(2q_2 + 1)}\right)\right) + \frac{2}{q_2^2} \left(1 + \log \left(\frac{31q_2^2 \log 999}{\log(2q_2^2 + 1)}\right)\right),$$

and this inequality does not hold when  $q_2 \geq 23$ .

#### 4.1 The conclusion of the proof of Theorem 2

Thus,  $3 \leq q_1 < q_2 \leq 19$ . The argument showing that  $\alpha_1 = 2$  except if  $(q_1, g) = (3, 729)$  now shows that  $\alpha_2 = 1$ . We are now able to show that  $s = 2$ . Indeed, if it were not so, then we would have both  $\tau(n/q_1) \geq 4$  and  $\tau(n/q_2) \geq 4$ . A quick computation with Mathematica shows that while there are pairs  $(q, g)$  such that  $\nu_q(gu_{q-1}) \geq 4$  in our ranges, there is no odd  $g$  in  $[3, 999]$  that has the above property with respect to two different primes  $3 \leq q_1 < q_2 \leq 19$ . Thus, either  $n = q_1 q_2$ , or  $n = 9q_2$  and  $g = 729$ . To test these last pairs, we proceeded as follows. First we have detected all pairs  $(n, g)$  with  $n = q_1 q_2$  with  $3 \leq q_1 < q_2 \leq 19$  and odd  $g \in [3, 999]$  such that  $\nu_{q_i}(gu_{n-1}) \geq 2$  holds for both  $i = 1, 2$ . There are 2043 such pairs. For each one of these we checked that  $\nu_2(u_{n-1}) < 14$ . Similarly, when  $Q_1 = 9$  and  $g = 729$ , the only possibility for  $q_2$  in our range such that  $\nu_{q_2}(u_{q_2-1}) \geq 2$  is  $q_2 = 11$ , but in this case  $n = 99$  and  $\nu_2(u_{n-1}) = 1 < 14$ . This finishes the proof of Theorem 2.

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