

# A NOTE ON PRODUCT SETS OF RATIONALS

JAVIER CILLERUELO

ABSTRACT. Bourgain, Konyagin and Shparlinski obtained a lower bound for the size of the product set  $AB$  when  $A$  and  $B$  are sets of positive rational numbers with numerator and denominator less or equal than  $Q$ . We extend and slightly improve that lower bound using a different approach.

## 1. INTRODUCTION

Bourgain, Konyagin and Shparlinsky [1] obtained a lower bound for the size of the product of two sets of rational numbers

$$A, B \subset \mathcal{F}_Q = \{q/q' : 1 \leq q, q' \leq Q\}$$

and they applied it to the study of the distribution of elements of multiplicative groups in residue rings. See [3] and [2] for related results and more applications of this useful inequality.

**Theorem A** (BKSh). *If  $A, B \subset \mathcal{F}_Q$  then*

$$(1) \quad |AB| \geq |A||B| \exp\left(- (9 + o(1)) \log Q / \sqrt{\log \log Q}\right),$$

where  $o(1) \rightarrow 0$  when  $Q \rightarrow \infty$ .

For any real numbers  $Q, Q' \geq 1$  let  $\mathcal{F}_{Q, Q'}$  denotes the set of rational numbers

$$\mathcal{F}_{Q, Q'} = \{q/q' : 1 \leq q \leq Q, 1 \leq q' \leq Q'\}.$$

We give the following result which extends and slightly improves Theorem A.

**Theorem 1.** *If  $A, B \subset \mathcal{F}_{Q, Q'}$  then*

$$|A/B| \geq |A||B| \exp\left(- (2\sqrt{\log 2} + o(1)) \log(QQ') / \sqrt{\log \log(QQ')}\right),$$

where  $o(1) \rightarrow 0$  when  $QQ' \rightarrow \infty$ .

Taking  $Q' = Q$  and the set  $1/B = \{b^{-1} : b \in B\}$  instead of  $B$  we improve the constant in (1).

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**Corollary 1.** *If  $A, B \in \mathcal{F}_Q$ , then*

$$|AB| \geq |A||B| \exp\left(-\left(4\sqrt{\log 2} + o(1)\right) \log Q / \sqrt{\log \log Q}\right).$$

## 2. PROOF OF THEOREM 1

For any pair of sets  $A, B \subset \mathcal{F}_{Q, Q'}$  and  $\gcd(r, s) = 1$  we define the sets

$$\begin{aligned} \mathcal{M}(A \times B, r/s) &= \{(a/a', b/b') \in A \times B : \gcd(a, b) = r, \gcd(a', b') = s\} \\ A_{r/s} &= \{a/a' \in A, r \mid a, s \mid a'\} \\ B_{r/s} &= \{b/b' \in B, r \mid b, s \mid b'\}. \end{aligned}$$

It is clear that  $\mathcal{M}(A \times B, r/s) \subset A_{r/s} \times B_{r/s}$ , so we have

$$(2) \quad |\mathcal{M}(A \times B, r/s)| \leq |A_{r/s}| |B_{r/s}|.$$

We claim that each  $c/d \in A/B$  (assume that  $\gcd(c, d) = 1$ ) has at most  $\tau(c)\tau(d)$  representation as

$$(3) \quad \frac{c}{d} = \frac{a/a'}{b/b'}$$

with  $(a/a', b/b') \in \mathcal{M}(A \times B, r/s)$ . Indeed we observe that (3) implies  $\frac{c}{d} = \frac{a_0 b'_0}{b_0 a'_0}$  where  $a_0 = a/r$ ,  $b_0 = b/r$ ,  $a'_0 = a_0/s$ ,  $b'_0 = b_0/s$ . Since  $\gcd(c, d) = 1$  and  $\gcd(a_0 b'_0, a'_0 b_0) = 1$  then  $c = a_0 b'_0$  and  $d = a'_0 b_0$ , which proves the claim.

Note that  $c = a_0 b'_0 \leq QQ'$  and  $d = a'_0 b_0 \leq QQ'$ , thus the claim implies the inequality

$$(4) \quad |\mathcal{M}(A, B, r/s)| \leq T^2 |A/B|,$$

where  $T = T(QQ')$  and  $T(x)$  is the function

$$T(x) = \max_{m \leq x} \tau(m).$$

Using (2), (4) and the well known inequality

$$\sum_{\substack{1 \leq r, s \\ rs \leq x}} 1 \leq x(1 + \log x)$$

we get

$$\begin{aligned} (5) \quad |A||B| &= \sum_{\substack{rs \leq x \\ (r, s) = 1}} |\mathcal{M}(A, B, r/s)| + \sum_{\substack{rs > x \\ (r, s) = 1}} |\mathcal{M}(A, B, r/s)| \\ &\leq T^2 |A/B| x(1 + \log x) + \sum_{\substack{rs > x \\ (r, s) = 1}} |A_{r/s}| |B_{r/s}| \end{aligned}$$

for any real number  $x \geq 1$ . If  $x$  is such that the last sum is less than  $|A||B|/2$  then we get

$$(6) \quad |A/B| \geq \frac{|A||B|}{2T^2x(1 + \log x)}.$$

Now we are ready to prove the key Lemma.

**Lemma 2.** *For any  $n \geq 1$  and for any  $A, B \in \mathcal{F}_{Q, Q'}$  with real numbers  $Q, Q' \geq 1$ , we have*

$$(7) \quad |A/B| \geq \frac{|A||B|}{(4T)^{n+1}(QQ')^{1/n}(1 + \log(QQ'))}$$

where  $T = \max_{m \leq QQ'} \tau(m)$ .

*Proof.* We proceed by induction on  $n$ : trivially, since  $|B| \leq QQ'$  we have

$$|A/B| \geq |A| \geq \frac{|A||B|}{QQ'},$$

which proves (7) for  $n = 1$ . Suppose that Lemma 2 is true for some  $n \geq 1$ .

If there is  $r/s$  such that

$$(8) \quad |A_{r/s}||B_{r/s}| \geq \frac{(QQ')^{\frac{1}{n(n+1)}}}{4T(rs)^{1/n}} |A||B|$$

we use induction for the sets  $A_{r/s}, B_{r/s} \in \mathcal{F}_{Q/r, Q'/s}$ . By observing that the function  $T(x) = \max_{m \leq x} \tau(m)$  is a non decreasing function we have

$$|A/B| \geq |A_{r/s}/B_{r/s}|$$

$$\begin{aligned} \text{(by induction hypothesis)} &\geq \frac{|A_{r/s}||B_{r/s}|}{(4T)^{n+1}((Q/r)(Q'/s))^{1/n}(1 + \log((Q/r)(Q'/s)))} \\ \text{(by (8))} &\geq \frac{|A||B|}{(4T)^{n+2}(QQ')^{1/(n+1)}(1 + \log(QQ'))}. \end{aligned}$$

Thus, we assume that

$$|A_{r/s}||B_{r/s}| < \frac{(QQ')^{\frac{1}{n(n+1)}}}{4T(rs)^{1/n}} |A||B|$$

for any  $r/s$ ,  $(r, s) = 1$ . In this case we have

$$\begin{aligned} \sum_{rs > x} |A_{r/s}||B_{r/s}| &\leq \max_{rs > x} (|A_{r/s}||B_{r/s}|)^{1/2} \sum_{rs > x} |A_{r/s}|^{1/2} |B_{r/s}|^{1/2} \\ (9) \quad &\leq \frac{(QQ')^{\frac{1}{2n(n+1)}}}{2T^{1/2}x^{\frac{1}{2n}}} (|A||B|)^{1/2} \left( \sum_{r,s} |A_{r/s}| \right)^{1/2} \left( \sum_{r,s} |B_{r/s}| \right)^{1/2}. \end{aligned}$$

To estimate the sums in the brackets we have

$$(10) \quad \sum_{r,s} |A_{r/s}| = \sum_{q/q' \in A} \sum_{\substack{r,s \\ r|q, s|q'}} 1 \leq \sum_{q/q' \in A} \tau(qq') \leq |A|T.$$

Putting in (9) the estimate (10) and the analogous for  $\sum_{r,s} |B_{r/s}|$  we have

$$\sum_{rs > x} |A_{r/s}| |B_{r/s}| \leq |A||B| \frac{T^{1/2} (QQ')^{\frac{1}{2n(n+1)}}}{2x^{\frac{1}{2n}}}.$$

Taking  $x = T^n (QQ')^{\frac{1}{n+1}}$  we get

$$\sum_{rs > x} |A_{r/s}| |B_{r/s}| \leq |A||B|/2.$$

Then (6) applies and noting that  $\log x \leq \log((QQ')^{n+\frac{1}{n+1}}) \leq 2n \log(QQ')$  we get

$$\begin{aligned} |A/B| &\geq \frac{|A||B|}{2T^2 x (1 + \log x)} \\ &\geq \frac{|A||B|}{2T^{n+2} (QQ')^{\frac{1}{(n+1)}} (1 + 2n \log(QQ'))} \\ &\geq \frac{|A||B|}{(4T)^{n+2} (QQ')^{\frac{1}{(n+1)}} (1 + \log(QQ'))} \times \frac{2^{2n+3} (1 + \log(QQ'))}{1 + 2n \log(QQ')} \\ &\geq \frac{|A||B|}{(4T)^{n+2} (QQ')^{\frac{1}{(n+1)}} (1 + \log(QQ'))}. \end{aligned}$$

□

The well known upper bound for the divisor function,

$$\tau(m) \leq \exp((\log 2 + o(1)) \log m / \log \log m)$$

implies

$$T \leq \exp((\log 2 + o(1)) \log(QQ') / \log \log(QQ')).$$

Thus, an optimal choice of  $n$  in Lemma 2 is  $n \sim \sqrt{\frac{\log \log(QQ')}{\log 2}}$ , from where Theorem 1 follows.

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INSTITUTO DE CIENCIAS MATEMÁTICAS (CSIC-UAM-UC3M-UCM) AND DEPARTAMENTO DE MATEMÁTICAS, UNIVERSIDAD AUTÓNOMA DE MADRID, 28049 MADRID, SPAIN